Domain Decomposition and Multilevel Methods in Diffpack

Are Magnus Bruaset, Hans Petter Langtangen, and Gerhard W. Zumbusch

1 Introduction

Looking back a decade or two, the computing power commonly available to scientists and engineers has grown at an amazing rate. Following this race for megaflops, the scientific community has developed a taste for solving mathematical problems of increasing levels of complexity. Naturally, this trend calls for sophisticated numerical methods that are capable of solving the problems in question in an efficient, yet reliable, way.

Since the bottleneck of many scientific applications turns out to be the numerical solution of linear or nonlinear systems of equations, this field has been subject to intensive research over the years. In particular, much attention has been paid to domain decomposition and multigrid methods, which have proven to be highly efficient strategies for many types of applications. Although the performance of such methods has been theoretically analyzed, extensive numerical experimentation is usually required in more complicated applications in order to obtain the best possible results. From this point of view, there is a need for software environments with genuine support for these types of numerical experiments, giving the user access to different algorithmic scenarios at the press of a button.

Domain decomposition and multilevel methods contain a variety of more standard numerical building blocks (linear solvers, matrix assembly, interpolation of fields, etc.). Successful software for complicated applications must offer the user a flexible run-time combination of all these different components. The purpose of the present paper is to describe how one can achieve such flexible software. In particular, we present a unified framework for domain decomposition and multilevel methods, and show how this framework can be efficiently implemented in existing software packages for PDEs.

The unified framework about to be presented is in part well known from the analysis

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1 The software design discussed in this paper has been implemented and verified using the object-oriented PDE library Diffpack [Dif, BL97].

Ninth International Conference on Domain Decomposition Methods
Editor Petter E. Bjørstad, Mauge S. Espedal and David E. Keyes ©1998 DDM.org
of overlapping and non-overlapping methods [DW90], as well as from theory for overlapping and multilevel schemes [Xu92]. In this context, the goal of this paper is to extend the known framework to cover even more methods in common use, especially some Schur complement and nonlinear schemes. We will formulate the framework in a novel way that encourages systematic implementation of a wide class of domain decomposition and multilevel methods. Finally, we report on the experiences gathered from a particular implementation in the Diffpack software.

2 The Unified Framework

*Abstract Schwarz Method*

We consider linear systems, \( Ax = f \), where \( A \) arises from the discretization of a partial differential operator. The prototype solution algorithm, the additive Schwarz method, can be written as

\[
B = \sum_j R_j^* B_j R_j
\]

with approximate solvers \( B_j \) operating on a subspace \( V_j \) and restrictions \( R_j \) and adjoint interpolations \( R_j^* \). In particular we consider the following methods and its variants (see also [SBG96]), which all can be written in the framework:

<table>
<thead>
<tr>
<th>multilevel iteration</th>
<th>additive, multiplicative, nonlinear</th>
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<tbody>
<tr>
<td>overlapping Schwarz</td>
<td>additive, multiplicative, nonlinear</td>
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<td></td>
<td>with, without coarse grid</td>
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<tr>
<td>Schur complement iteration</td>
<td>exact, inexact local solves</td>
</tr>
<tr>
<td>Schur complement preconditioner</td>
<td>Neumann-Neumann, wirebasket</td>
</tr>
<tr>
<td></td>
<td>with, without coarse grid</td>
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The interface basically consists of

- transfer or restriction operators such as \( R_j \) and \( R_j^* \),
- (approximate) subspace solvers \( B_j \),
- evaluation of residuals \( f - A_j x \) on a subdomain (optional).

Here, \( A_j \) is the discrete operator on \( V_j \). The appropriate solver \( B_j \) is normally chosen as a traditional linear or nonlinear method, whereas the “transfer” \( R_j \) is usually implemented via sparse matrices or difference stencils. However, also algorithmic implementations of \( R_j \) and message passing in a parallel environment are possible. Software components for \( B_j \), \( R_j \), \( R_j^* \) and \( f - A_j x \) are normally found in a package for PDEs.

*Multilevel Methods*

Standard multigrid algorithms fit into the frame outlined in 2. In the linear multigrid context the abstraction of the (not necessarily adjoint) grid transfer \( R_j \), \( R_j^* \) and the local pre- and post-smoothers \( B_j \), \( B_j^* \) is sometimes referred to as abstract multigrid.
The implementation follows equation (4.1.1) in Hackbusch [Hac85]. Using multigrid as a preconditioner implies the initial guess \( x = 0 \).

However, it is also possible to include nonlinear multigrid with nonlinear smoothers \( B_j, B_j^* \) [Zum96a]. Using equation (9.3.3) in [Hac85], we can implement the nonlinear FAS scheme and the nested nonlinear multigrid version by Hackbusch. The pre- and post-smoothers are now nonlinear iterative solvers.

This abstract multilevel approach can be used for different variants of multigrid. The particular algorithm depends on the initialization and implementation of the smoothers \( B_j \), the transfer operators \( R_j \), and the operators \( A_j \). There may be non-nested, non-matching, or adaptively refined grids, operators defined by Galerkin products or operator dependent transfers, or algebraic multigrid transfers and operators.

**Schur Complement Methods**

We decompose the stiffness matrix into one part related to the coupling interface \( c \) and several independent parts related to the interior nodes of the subdomains \( j = 1, \ldots, m \).

\[
A = \begin{pmatrix}
A_{11} & A_{1c} \\
A_{c1} & A_{c2} \\
\vdots & \vdots \\
A_{c1} & A_{c2} \\
\end{pmatrix}
\]

The Schur complement is defined by \( S = A_{cc} - \sum_j A_{cj} A_{jj}^{-1} A_{jc} \).

In the case we solve subdomain problems exactly or accurate enough, or we are able to compute \( S \) itself rather than the action of \( S \), the equation system reduces to a system \( S x_c = f_c \) for the unknowns on the interface \( x_c \).

**Neumann-Neumann Preconditioner for the Schur Complement.** Preconditioning of the interface system with a Neumann-Neumann algorithm [BGLV89] looks like equation 2.1 in the framework of additive Schwarz methods. The operators \( B_j \) now solve homogeneous Neumann problems obtained by sub-assembly on a subdomain. The transfer operators \( R_j \) and \( R_j^* \) copy (scaled) nodal data from the interface \( x_c \) to the nodes on the boundary of the subdomain and vice versa. Hence the ordinary additive Schwarz implementation can be used.

A coarse grid can be added to the Neumann-Neumann preconditioner without changes of the algorithm [DW95] just like in the overlapping Schwarz case adding one subspace \( j = 0 \). Coupling a coarse grid equation in a multiplicative symmetric way instead, called balancing [Man93], requires some extensions resulting in a mixture of the additive and the multiplicative Schwarz algorithm. One can use the standard global-to-local communication pattern and additional residual evaluations to create the multiplicative coupling.

**Wire Basket Preconditioner for the Schur Complement.** A preconditioner of wirebasket type [Smi91] for the Schur complement may also be written as an additive Schwarz method now for subspaces of the interface consisting of vertices \( v \), single faces \( f \), and single edges \( e \),

\[
B_c = R^v \star B^v R^v + \sum_j R^e_j \star B^e_j R^e_j + \sum_k R^f_k \star B^f_k R^f_k
\]
The transfer operators $R^e_j$ and $R^f_j$ perform a hierarchical basis transform. The vertex solver $B^e$ can be implemented by a standard coarse grid solver and the edge solvers $B^e_j$ can be substituted by diagonal scaling. However, the face solvers $B^f_j$ have to be implemented as preconditioners for an interface problem covering both adjacent subdomains.

The two-dimensional analog BPS has a similar structure. It can be implemented by leaving out the faces and choosing some local interface preconditioner for the edge solvers $B^e_j$ \cite{BPS86}.

**Inexact Schur Complement.** Introducing approximate local Dirichlet solvers $B_j$ in the computation of the Schur complement $S \approx A_{cc} - \sum_j A_{cj} B_j A_{jc}$, one cannot restrict the computations to $S x_c = f_c$, but the full system $A x = f$ must be considered. We seek a preconditioner for the matrix $A$, which is constructed using a preconditioner $B_e$ for $S$ and preconditioners $B_j$ for the local Dirichlet type problems $A_{jj}$ that may differ from $B_j$ used for the Schur complement. However, the inexact solvers $B_j$ in the Schur complement method are still not fully understood theoretically \cite{HLM91}.

We write the Schur decomposition in the additive Schwarz framework

$$B = R^e_c B_e R_e + \sum_j R^e_j B_j R_j$$

with restrictions $R_j x = x_j$ and the transformation to the Schur system $R_c x = x_c - \sum_j A_{cj} B_j x_j$.

One can use multilevel and domain decomposition methods or standard iterative solvers for the implementation of the local solvers $B_j$, $B^*_j$, $B\hat j_j$, and the local solvers used in $B_e$. This approach to the Schur complement can also be used for nonlinear problems in a Picard iteration manner (see \cite{Zum96b}) since the residual for the full system has to be evaluated every outer iteration step anyway. An exact Dirichlet solver $B_j$ leads to zero residuals at interior nodes and is more efficiently implemented computing on the interface directly as in the previous section.

3 Realizing the Framework

As mentioned initially, the abstract view of the different domain decomposition and multilevel methods taken in the description above has been realized in terms of software. More precisely, the described framework is the basis for the implementation of domain decomposition and multilevel algorithms in the Diffpack simulation environment. Diffpack \cite{Dif} consists of a set of object-oriented libraries written in C++, intended to simplify the implementation of PDE solvers \cite{BL97}. The involved libraries contain many useful abstractions that come quite natural to developers of PDE software. For instance, the application programmer has immediate access to high-level abstractions for linear systems, various matrix formats, different solvers and preconditioners, as well as building blocks for finite element and finite difference discretizations.

Originally, Diffpack was designed without attention to domain decomposition or multilevel algorithms. Nevertheless, by adopting a unified approach to such methods, one can in principle implement this functionality as an add-on module to existing PDE packages. However, it is then crucial that the framework organizing the different
methods can take advantage of tools already present in the available software platform. During the course of the multilevel extension of Diffpack, we have experienced that a clean design of such an add-on module depends heavily on a clean design and modular structure of the underlying libraries. We believe that abstract data types and object-oriented programming are important mechanisms for achieving the necessary degree of modularity. In fact, the Diffpack code for the described framework was realized as a high-level, compact combination of existing C++ classes. We refer to [ABL97, BL97, ABC+97, BL96] and the references therein for information about object-oriented numerics, the efficiency of C++ for scientific computing, the design of Diffpack and examples on Diffpack applications. It should be mentioned that object-oriented implementations of domain decomposition and multilevel strategies have also been addressed by other authors, e.g. as part of the PETSc system [GS94, PET].

Equation solvers in Diffpack utilize abstractions for linear systems, as well as linear and nonlinear solvers. In this context, a linear system consists of a coefficient matrix, a solution vector, a right-hand side and possibly a preconditioner, which will be reused in our framework. The preconditioner can either be a matrix or an action. In case of an action, the preconditioner can, e.g., call a linear solver for the same or a related PDE problem. Linear solvers can utilize convergence monitors in order to control the degree of solution accuracy [BL96]. At its present stage of development, Diffpack offers relatively simple nonlinear solvers, like Newton's method or the Picard iteration. This type of algorithm requires the programmer to define and solve a linear (sub)system. Operators can be defined in terms of coefficient matrices arising from finite element assembly. Since a finite element grid is just a C++ object in Diffpack, it is easy to create a hierarchy of grids, and apply toolboxes for finite element schemes and the PDE's definition to create the associated operators and right-hand sides.

The layered, modular design of Diffpack building blocks can be immediately applied to create linear and nonlinear operators, smoothers, transfer operators, residuals and other basic components needed in domain decomposition and multilevel methods, reusing existing code. For instance, the subspace solver $B_j$ in a domain decomposition method may be implemented as a linear solver call or a nonlinear solution procedure, both present in the original package. To allow maximum flexibility, the programmer of the PDE application is responsible for defining $B_j$, and contrary to the PETSc approach [GS94, PET], we do not require a reference to a linear solver object. The main ingredients of the domain decomposition interface are the procedures for grid transfer and for local solvers. Data is only stored for some auxiliary vectors and parameters like the multigrid cycle type ($V, W, \ldots$).

The technical details of taking an existing Diffpack application and equipping it with domain decomposition and multilevel methods are described elsewhere [Zum96a, Zum96b]. However, due to the flexibility of the original software components, it turns out to be trivial to run a multigrid solver and experiment with various pre- and post-smoothers (choice of algorithm, number of sweeps, order of unknowns), coarse grid solvers (iterative and direct, grid size), cycle-types, nested iterations, non-matching grids, semi-coarsening, multigrid used as a preconditioner or as a stand-alone solver, different nonlinear versions, grid types and special procedures to initialize operators. For domain decomposition, the type, precision and termination of subdomain solvers, the decomposition of the domain, the type of a coarse grid and coarse-grid solver, and the scaling of transfer operators are of main
interest. Consequently, the resulting software environment satisfies the most important requirement stated in the introduction of this paper; to offer the user genuine support for systematic numerical experiments with sophisticated multilevel strategies.

As previously mentioned, Schur complement methods do not immediately fit into a unified framework. This is also reflected in the pilot implementation. We use an implicit representation of the Schur complement, implemented as an algorithmically defined matrix. This is basically a new type of matrix in Diffpack, realized as a subclass in the existing matrix hierarchy [BL96], which implements the matrix product $Sx$ only. New and old application software can of course work with this matrix type through an abstract (base class) interface. The action $Sx$ is implemented by calling (reusing) subdomain Dirichlet solvers for $A^{-1}_{jj}$ and several matrix multiplications, since the concept of a Schur complement was not present in Diffpack. Direct access to $S$ is then not available. However, standard Schwarz preconditioners can be used for the Schur complement system, reusing even the domain decomposition algorithms.

The domain decomposition and multilevel methods introduced in Diffpack are organized in a class hierarchy with \texttt{DDSolver} as base class, see Figure 1. Various specific solution strategies are organized as subclasses of either (multiplicative) \texttt{Multigrid} or (alternating) Schwarz domain decomposition (\texttt{SchwarzDD}). \texttt{Multigrid} uses level to level data transfer ($j \rightarrow j+1$), while the \texttt{SchwarzDD} algorithms uses local to global data transfer. \texttt{SymSchwarzDD} is a symmetric version of the multiplicative \texttt{SchwarzDD}, cycling back and forth like the multiplicative \texttt{Multigrid}. The \texttt{CoarseAddCoarseSchwarzDD} version is used for special treatment of the coarse grid, while the subdomains are treated as in the additive \texttt{AddSchwarzDD}. The hierarchy itself is designed to optimize code reuse and reflects software issues, while one can certainly think of modifications in the categorization. The subclasses make use of existing solvers and preconditioners in Diffpack, while still offering the duality of being accessible as new solvers and new preconditioners in the original libraries. For the authors, this experience of playing

Figure 1 Multilevel and Domain Decomposition algorithms. Abstract \texttt{DDSolver}, the multiplicative and additive multigrid algorithms including nonlinear versions and additive, multiplicative, symmetric multiplicative and mixed additive/multiplicative Schwarz algorithms.
around with abstractions and extending libraries in ways that were not initially planned for, has been a strong indication that modern programming techniques, such as object-oriented programming, are vital for an accelerated development of scientific computing.

4 Efficiency

We have outlined a flexible software framework for domain decomposition and multilevel methods, where the particular Diffpack implementation is in C++. Many will expect the computational efficiency of such flexible implementations and the use of C++ to be significantly worse than special-purpose Fortran codes tailored at a specific PDE and solution algorithm. To shed some light on this problem we have performed some simple numerical experiments with multigrid methods for a Poisson equation, with smooth variable coefficients, on the unit square. A general Diffpack implementation, also applicable to unstructured grids, was compared to (a) the adaptive PLTMG code [Bau94], (b) a finite difference based example Fortran code with MPI [Dou95a], and (c) a constant 5-point stencil sparse matrix Fortran Poisson solver Madpack5 [Dou95b]. The table below shows CPU times for various problem sizes.

<table>
<thead>
<tr>
<th>level $j$</th>
</tr>
</thead>
<tbody>
<tr>
<td>size $n$</td>
</tr>
<tr>
<td>4</td>
</tr>
<tr>
<td>PLTMG</td>
</tr>
<tr>
<td>MPI example</td>
</tr>
<tr>
<td>madpack5</td>
</tr>
<tr>
<td>Diffpack</td>
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</tbody>
</table>

As we see, there is no indication that the object-oriented implementation style in C++ implies a significant loss of computational efficiency. The reasons for this are simple; object-orientation is only used for high-level administration in Diffpack, whereas CPU-time consuming operations usually take place in low level C/Fortran-style routines that can be highly optimized by today’s compiler technology. Moreover, the general finite element software in Diffpack makes use of simplified, optimized algorithms when it is known that the grid is a uniform lattice.

5 Conclusion

We have outlined a unified framework for the whole set of domain decomposition and multilevel methods. The framework has been realized in Diffpack using object-oriented programming techniques. We have indicated that the implementation has the same level of efficiency as tailored, traditional implementations in Fortran or C, but with much more flexibility and extensibility.
REFERENCES


