Fixing Numbers of Graphs and Groups

Courtney R. Gibbons

University of Nebraska – Lincoln Department of Mathematics 228 Avery Hall PO Box 880130 Lincoln, NE 68588-0130 Joshua D. Laison Mathematics Department Willamette University 900 State St. Salem, OR 97301

jlaison@willamette.edu

s-cgibbon5@math.unl.edu

Submitted: Sep 11, 2006; Accepted: Mar 12, 2009; Published: Mar 20, 2009 Mathematics Subject Classification: 05C25

Abstract

The fixing number of a graph G is the smallest cardinality of a set of vertices S such that only the trivial automorphism of G fixes every vertex in S. The fixing set of a group Γ is the set of all fixing numbers of finite graphs with automorphism group Γ . Several authors have studied the distinguishing number of a graph, the smallest number of labels needed to label G so that the automorphism group of the labeled graph is trivial. The fixing number can be thought of as a variation of the distinguishing number in which every label may be used only once, and not every vertex need be labeled. We characterize the fixing sets of finite abelian groups, and investigate the fixing sets of symmetric groups.

1 Introduction

In this paper we investigate breaking the symmetries of a finite graph G by labeling its vertices. There are two standard techniques to do this. The first is to label all of the vertices of G with k distinct labels. A labeling is **distinguishing** if no nontrivial automorphism of G preserves the vertex labels. The **distinguishing number** of G is the minimum number of labels used in any distinguishing labeling [1, 13]. The **distinguishing chromatic number** of G is the minimum number of labels used in any distinguishing labeling which is also a proper coloring of G [6].

The second technique is to label a subset of k vertices of G with k distinct labels. The remaining labels can be thought of as having the null label. We say that a labeling of G is *fixing* if no non-trivial automorphism of G preserves the vertex labels, and the *fixing* number of G is the minimum number of labels used in any fixing labeling.

2 Fixing Graphs

More formally, suppose that G is a finite graph and v is a vertex of G. The **stabilizer** of v, stab(v), is the set of group elements $\{g \in \operatorname{Aut}(G) | g(v) = v\}$. The **(vertex) stabilizer** of a set of vertices $S \subseteq V(G)$ is stab $(S) = \{g \in \operatorname{Aut}(G) | g(v) = v \text{ for all } v \in S\}$. A vertex v is **fixed** by a group element $g \in \operatorname{Aut}(G)$ if $g \in \operatorname{stab}(v)$. A set of vertices $S \subseteq V(G)$ is a **fixing set** of G if stab(S) is trivial. In this case we say that S **fixes** G. The **fixing number** fix(G) of a graph G is the smallest cardinality of a fixing set of G [3, 5, 9].

Equivalently, S is a fixing set of the graph G if whenever $g \in Aut(G)$ fixes every vertex in S, g is the identity automorphism. A set of vertices S is a **determining set** of G if whenever two automorphisms $g, h \in Aut(G)$ agree on S, then they agree on G, i.e., they are the same automorphism [3]. The following lemma shows that these two definitions are equivalent.

Lemma 1. A set of vertices is a fixing set if and only if it is a determining set.

Proof. Suppose that S is a determining set. Since the identity automorphism e fixes every vertex in S, then by the definition of a determining set, every other element $g \in \operatorname{Aut}(G)$ that fixes every vertex in S must be the identity. Therefore S is a fixing set. Conversely, suppose that S is a fixing set. Let g and h agree on S. Then $g^{-1}h$ must fix every element in S. Hence by the definition of a fixing set, $g^{-1}h = e$, so g = h. Therefore S is a determining set.

Suppose G is a graph with n vertices. Since fixing all but one vertex of G necessarily fixes the remaining vertex, we must have $fix(G) \leq n-1$. In fact, suppose that any n-2 vertices have been fixed in G, yet G still has a non-trivial automorphism. Then this automorphism must be the transposition of the remaining two vertices. This implies that the only graphs which have fix(G) = n-1 are the complete graphs and the empty graphs. On the other hand, the graphs with fix(G) = 0 are the *rigid graphs* [1], which have trivial automorphism group. In fact, almost all graphs are rigid [2], so most graphs have fixing number 0.

The **orbit** of a vertex v, $\operatorname{orb}(v)$, is defined to be the set of vertices $\{w \in V(G) \mid g(v) = w \text{ for some } g \in \operatorname{Aut}(G)\}$. The Orbit-Stabilizer Theorem says that for any vertex v in G, $|\operatorname{Aut}(G)| = |\operatorname{stab}(v)| |\operatorname{orb}(v)|$ [12]. So when we are building a minimal fixing set of G, heuristically it makes sense to choose vertices with orbits as large as possible. This leads us to consider the following algorithm for determining the fixing number of a finite graph G:

The Greedy Fixing Algorithm.

- 1. Find a vertex $v \in G$ with $|\operatorname{stab}(v)|$ as small as possible (equivalently, with $|\operatorname{orb}(v)|$ as large as possible).
- 2. Fix v and repeat.
- 3. Stop when the stabilizer of the fixed vertices is trivial.

The set of vertices fixed by the greedy fixing algorithm must be a fixing set. We define the **greedy fixing number** $\operatorname{fix}_{greedy}(G)$ of the graph G to be the number of vertices fixed by the greedy fixing algorithm.

Open Question. Is $fix_{greedy}(G)$ well-defined for every finite graph G? In other words, is there a finite graph for which two different choices in Step 1 of the greedy fixing algorithm produce two different fixing sets of different sizes?

If $fix_{greedy}(G)$ is well-defined, we must have $fix(G) \leq fix_{greedy}(G)$. We use this same technique to derive upper bounds on the fixing sets of groups in the next section.

Open Question. Assuming $fix_{greedy}(G)$ is well-defined, is there a graph G for which $fix(G) \neq fix_{greedy}(G)$?

3 Fixing Sets of Groups

Following Albertson and Collins' exposition of distinguishing sets of groups [1], we define the **fixing set** of a finite group Γ to be $\operatorname{fix}(\Gamma) = {\operatorname{fix}(G) | G \text{ is a finite graph with } \operatorname{Aut}(G) \cong \Gamma}$. Our goal for the remainder of the paper is to find the fixing sets of a few well-known finite groups. We begin by describing two procedures that can be used to generate specific examples.

For every graph G, the natural representation of the elements of $\operatorname{Aut}(G)$ as permutations of the vertices of G is a group action of the group $\operatorname{Aut}(G)$ on the set V(G). Furthermore, $\operatorname{Aut}(G)$ acts **faithfully** on G, i.e., the only element of $\operatorname{Aut}(G)$ that fixes every vertex in G is the identity element. A group action of Γ on a graph G is **vertextransitive** if, given any two vertices $u, v \in V(G)$, there is an element of Γ that sends uto v. The following theorem appears in [7].

Theorem 2. Let Γ be a finite group. The set of vertex-transitive actions of Γ on all possible sets of vertices V is in one-to-one correspondence with the conjugacy classes of subgroups of Γ . Specifically, if v is any vertex in V, the action of Γ on V is determined by the conjugacy class of stab(v).

Suppose that Γ is the automorphism group of a graph G. Then Γ acts transitively on each orbit of the vertices of G under Γ . Hence given a group Γ , to find a graph Gwith automorphism group Γ , we choose a set of subgroups of Γ and generate the orbits of vertices of G corresponding to these subgroups using Theorem 2. There are two aspects of this construction which make the procedure difficult. First, the action of Γ on the entire graph G must be faithful for Γ to be a valid automorphism group. Second, after we construct orbits of vertices, we must construct the edges of G so that the set of permutations of vertices in Γ is exactly the set of edge-preserving permutations of G. However, this is not always possible.

An alternative approach uses the Orbit-Stabilizer Theorem. Given a graph G and a fixing set S of G, we order the elements of S as, say, v_1, \ldots, v_k , and we consider the chain of subgroups $e = \operatorname{stab}(\{v_1, \ldots, v_k\}) \leq \operatorname{stab}(\{v_1, \ldots, v_{k-1}\}) \leq \ldots \leq \operatorname{stab}(v_1) \leq \operatorname{Aut}(G)$.

If $o(v_i)$ is the number of vertices in $\operatorname{orb}(v_i)$ under the action of $\operatorname{stab}(\{v_1, \ldots, v_{i-1}\})$, then $|\operatorname{stab}(\{v_1, \ldots, v_{i-1}\})| = o(v_i)|\operatorname{stab}(\{v_1, \ldots, v_i\})|$. So $|\operatorname{Aut}(G)| = \prod_{1 \le i \le k} o(v_i)$. Hence given a finite group Γ , to find a graph G with automorphism group Γ and fixing number k, we choose a sequence of orbit sizes $(o(v_1), \ldots, o(v_k))$ whose product is $|\Gamma|$ and look for a graph with these orbit sizes. Both of these procedures were used to generate examples given below.

We now prove a few theorems valid for the fixing set of any finite group. Let Γ be a group generated by the set of elements $\mathcal{G} = \{g_1, g_2, \dots, g_k\}$. The **Cayley graph** $C(\Gamma, \mathcal{G})$ of Γ with respect to the generating set \mathcal{G} is a directed, edge-labeled multigraph with a vertex for each element of Γ , and a directed edge from the group element h_1 to the group element h_2 labeled with the generator $g \in \mathcal{G}$ if and only if $gh_1 = h_2$.

We obtain an undirected, edge-unlabeled graph $F(\Gamma, \mathcal{G})$ from the Cayley graph $C(\Gamma, \mathcal{G})$ by replacing each directed, labeled edge of $C(\Gamma, \mathcal{G})$ with a "graph gadget" so that $F(\Gamma, \mathcal{G})$ has the same automorphisms as $C(\Gamma, \mathcal{G})$. This technique is due to Frucht [10, 11] and is outlined in greater detail in [2]. An example is shown in Figure 1. We call $F(\Gamma, \mathcal{G})$ the **Frucht Graph** of Γ with respect to the generating set \mathcal{G} . The following lemma is easy to prove and also follows from the exposition in [2].

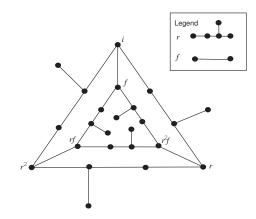


Figure 1: The Frucht graph $F(D_3, \{r, f\})$.

Lemma 3. For any group Γ and any generating set \mathcal{G} of Γ , $\operatorname{Aut}(C(\Gamma, \mathcal{G})) = \Gamma$ and $\operatorname{Aut}(F(\Gamma, \mathcal{G})) = \Gamma$. Furthermore, for two elements $g, h \in \Gamma$, the automorphism g takes the vertex h to the vertex gh in both $C(\Gamma, \mathcal{G})$ and $F(\Gamma, \mathcal{G})$.

Corollary 4. If G is a Cayley graph or a Frucht graph of a non-trivial group, then fix(G) = 1.

Proof. Suppose $G = F(\Gamma, \mathcal{G})$ for some group Γ (the argument for Cayley graphs is completely analogous). Since $\operatorname{Aut}(G) = \Gamma$ by Lemma 3, and Γ is not trivial by hypothesis, $\operatorname{fix}(G) > 0$. Now let h be an element of Γ (and so also a vertex in G). For any non-identity element $g \in \Gamma$, by Lemma 3, $g(h) = gh \neq h$. Thus $\operatorname{stab}(h)$ is trivial, and the single-vertex set $\{h\}$ is a fixing set of G.

In fact, the proof of Corollary 4 implies that every vertex of a Cayley graph is a fixing set, and every non-gadget vertex of a Frucht graph is a fixing set.

Corollary 5. For any non-trivial finite group Γ , $1 \in fix(\Gamma)$.

The *length* $l(\Gamma)$ of a finite group Γ is the maximum number of subgroups in a chain of subgroups $e < \Gamma_1 < \Gamma_2 < \ldots < \Gamma_{l(\Gamma)} = \Gamma$ [4].

Proposition 6. For any finite group, $\max(\operatorname{fix}(\Gamma)) \leq l(\Gamma)$.

Proof. If Γ is trivial, it has length 0 and fixing set $\{0\}$. Now suppose Γ is non-trivial, and let G be a graph with $\operatorname{Aut}(G) = \Gamma$. We fix a vertex v_1 in G with orbit larger than one. By the Orbit-Stabilizer Theorem, $\operatorname{stab}(v_1)$ is a proper subgroup of Γ . If we can find a different vertex v_2 with orbit greater than one under the action of $\operatorname{stab}(v_1)$, we fix v_2 . We continue in this way until we have fixed G. Since at each stage, $\operatorname{stab}(\{v_1, \ldots, v_i\})$ is a proper subgroup of $\operatorname{stab}(\{v_1, \ldots, v_{i-1}\})$, we cannot have fixed more than the length of the group.

Corollary 7. Let k be the number of primes in the prime factorization of $|\Gamma|$, counting multiplicities. Then $\max(\operatorname{fix}(\Gamma)) \leq k$.

Example 8. The graph C_6 has automorphism group D_6 and fixing number 2. The graph $C_3 \cup P_2$ has automorphism group D_6 and fixing number 3. On the other hand, $|D_6| = 12 = 2 \cdot 2 \cdot 3$. Hence fix $(D_6) = \{1, 2, 3\}$ by Corollaries 5 and 7.

Example 9. The graph shown in Figure 2 has automorphism group A_4 and fixing number 2. On the other hand, $|A_4| = 12 = 2 \cdot 2 \cdot 3$. So $\{1,2\} \subseteq \text{fix}(A_4) \subseteq \{1,2,3\}$, again by Corollaries 5 and 7. Lemma 10 shows that $3 \notin \text{fix}(A_4)$, so in fact fix $(A_4) = \{1,2\}$.

Lemma 10. There is no graph G with fix(G) = 3 and $Aut(G) = A_4$.

Proof. Suppose by way of contradiction that G is a graph with fix(G) = 3 and $Aut(G) = A_4$. Let $S = \{v_1, v_2, v_3\}$ be a minimum size fixing set of G. Note that $stab(v_1)$, $stab(v_2)$, and $stab(v_3)$ are all proper subgroups of A_4 . Therefore they must be isomorphic to \mathbb{Z}_2 , $\mathbb{Z}_2 \times \mathbb{Z}_2$, or \mathbb{Z}_3 . But if any of them have order less than 4, fixing that vertex and one other will fix G, and fix(G) = 2. So $stab(v_1) \cong stab(v_2) \cong stab(v_3) \cong \mathbb{Z}_2 \times \mathbb{Z}_2$. But there is only one copy of $\mathbb{Z}_2 \times \mathbb{Z}_2$ in A_4 , so $stab(v_1) = stab(v_2) = stab(v_3)$, and this subgroup must therefore also equal $stab(\{v_1, v_2, v_3\})$. So $\{v_1, v_2, v_3\}$ is not a fixing set of G, which is a contradiction.

Lemma 11. Suppose G is a graph, $\Gamma = \operatorname{Aut}(G)$ is a finite non-trivial group, and $g \in \Gamma$ is an element of order p^k , for p prime and k a positive integer. Then there exists a set of p^k vertices v_1, \ldots, v_{p^k} in G such that, as a permutation of the vertices of G, g contains the cycle $(v_1 \ldots v_{p^k})$.

Proof. Since g has order p^k , the cycle decomposition of g must include a cycle of length p^k . Label these vertices v_1, \ldots, v_{p^k} .

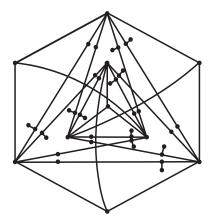


Figure 2: A graph G with $Aut(G) = A_4$ and fix(G) = 2.

Recall that the cartesian product of two groups Γ_1 and Γ_2 is the group $\Gamma_1 \times \Gamma_2 = \{(g,h) \mid g \in \Gamma_1, h \in \Gamma_2\}$ with group operation defined by $(g_1,h_1)(g_2,h_2) = (g_1g_2,h_1h_2)$. Recall also that the sum of two sets S and T is $S + T = \{s + t \mid s \in S, t \in T\}$.

Lemma 12. If Γ_1 and Γ_2 are finite non-trivial groups, then $fix(\Gamma_1) + fix(\Gamma_2) \subseteq fix(\Gamma_1 \times \Gamma_2)$.

Proof. Let $a \in \text{fix}(\Gamma_1)$ and $b \in \text{fix}(\Gamma_2)$. Then there exist graphs G_1 and G_2 with $\text{Aut}(G_1) = \Gamma_1$, $\text{Aut}(G_2) = \Gamma_2$, $\text{fix}(G_1) = a$, and $\text{fix}(G_2) = b$. Let G'_2 be the graph obtained from G_2 by attaching the graph Y_k shown in Figure 3 for some large value of k (for example, $|G_1| + |G_2|$) to each vertex of G_2 at the vertex a. Now consider the graph $H = G_1 \cup G'_2$, the disjoint union of the graphs G_1 and G'_2 . This graph has no automorphisms that exchange vertices between G_1 and G_2 , so we must have $\text{Aut}(H) \cong \text{Aut}(G_1) \times \text{Aut}(G'_2) \cong \text{Aut}(G_1) \times \text{Aut}(G_2) \cong \Gamma_1 \times \Gamma_2$. Furthermore, H is fixed if and only if both G_1 and G_2 are fixed, so fix(H) = a + b. Therefore $a + b \in \text{fix}(\Gamma_1 \times \Gamma_2)$.



Figure 3: The graph Y_k in the proof of Lemma 12 is shown on the left, and the graph A_k in the proof of Theorem 14 is shown on the right.

Note that for two finite non-trivial groups Γ_1 and Γ_2 , $1 \in fix(\Gamma_1 \times \Gamma_2)$ but $1 \notin fix(\Gamma_1) + fix(\Gamma_2)$.

Open Question. Is it true that for all finite non-trivial groups Γ_1 and Γ_2 , fix (Γ_1) + fix $(\Gamma_2) =$ fix $(\Gamma_1 \times \Gamma_2) \setminus \{1\}$?

3.1 Abelian groups

Lemma 13. If p is prime and k is a positive integer, then $fix(\mathbb{Z}_{p^k}) = \{1\}$.

Proof. By Corollary 5, $1 \in \text{fix}(\mathbb{Z}_{p^k})$. Conversely, suppose that there exists a graph G such that $\text{Aut}(G) = \mathbb{Z}_{p^k}$. By Lemma 11, there exists a vertex in G with orbit size p^k . By the Orbit-Stabilizer Theorem, fixing this vertex must fix the graph.

Let Γ be a finite abelian group with order n, and let $n = p_1^{i_1} \cdots p_k^{i_k}$ be the prime factorization of n. Recall that there is a unique factorization $\Gamma = \Lambda_1 \times \cdots \times \Lambda_k$, where $|\Lambda_j| = p_j^{i_j}, \Lambda_j = \mathbb{Z}_{p_j^{\alpha_1}} \times \cdots \times \mathbb{Z}_{p_j^{\alpha_t}}$, and $\alpha_1 + \ldots + \alpha_t = i_j$. The numbers $p_j^{\alpha_r}$ are called the *elementary divisors* of Γ [8].

Theorem 14. Let Γ be a finite abelian group, and let k be the number of elementary divisors of Γ . Then fix $(\Gamma) = \{1, \ldots, k\}$.

Proof. Let $\Gamma = \Gamma_1 \times \ldots \times \Gamma_k$ be the elementary divisor decomposition of Γ . For every $1 \leq i \leq k$, let $H_i = F(\Gamma_i \times \ldots \times \Gamma_k, \mathcal{G})$ be any Frucht graph of $\Gamma_i \times \ldots \times \Gamma_k$. There are an infinite number of finite graphs with automorphism group \mathbb{Z}_n and fixing number 1; for example, every graph in the family of graphs shown in Figure 4 has automorphism group \mathbb{Z}_5 and fixing number 1. We may therefore let G_1, \ldots, G_k be distinct graphs, not isomorphic to H_i for any i, with automorphism groups $\Gamma_1, \ldots, \Gamma_k$, respectively, and fixing number 1. Let G be the disjoint union $(\bigcup_{j=1}^{i-1} G_j) \cup H_i$. We also choose G_1, \ldots, G_k so that no automorphism of G moves a vertex from one G_j to another, or from any G_j to H_i , or vice versa. The graphs shown in Figure 4 are examples of graphs G_j which have this property.

Then G has automorphism group Γ . Furthermore, every fixing set of G must include at least one vertex from each subgraph G_j and at least one vertex from H_i , and any set with exactly one vertex moved by an automorphism from each G_j and from H_i is a fixing set of G. Therefore fix(G) = i. Since we have constructed a graph G with $\operatorname{Aut}(G) = \Gamma$ and fix(G) = i for any $1 \leq i \leq k$, $\{1, \ldots, k\} \subseteq \operatorname{fix}(\Gamma)$.

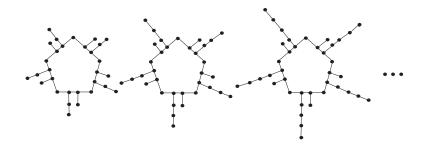


Figure 4: An infinite family of graphs with automorphism group \mathbb{Z}_5 and fixing number 1.

We prove the reverse inclusion by induction. Suppose Γ is a finite abelian group and G is a finite graph with $\operatorname{Aut}(G) = \Gamma$. If Γ has one elementary divisor, then the result follows from Lemma 13. Suppose that Γ has k > 1 elementary divisors. We choose an

elementary divisor p^m of Γ . Then $\Gamma = \mathbb{Z}_{p^m} \times \Gamma'$ for a smaller finite abelian group Γ' . Let g be a generator of the subgroup \mathbb{Z}_{p^m} of Γ . By Lemma 11, there exists a set of p^m vertices v_1, \ldots, v_{p^m} in G such that, as a permutation of the vertices of G, g contains the cycle $(v_1 \ldots v_{p^m})$.

Let H be the connected component of G containing v_1 . If H is a tree, let G' be the graph obtained from G by attaching the graph $A_{|G|}$ shown in Figure 3 to G by identifying the vertex a in $A_{|G|}$ with the vertex v_1 in G. Otherwise, let G' be the graph obtained from G by attaching the graph $Y_{|G|}$ shown in Figure 3 to G by identifying the vertex a in $Y_{|G|}$ with the vertex v_1 in G. Denote the subgraph $A_{|G|}$ or $Y_{|G|}$ in G' by H'. We claim that $\operatorname{Aut}(G')$ is a subgroup of Γ' . First, we show that G' does not have any additional automorphisms that G does not have. Suppose h is an automorphism of G' and not G. So h must move some vertex of H'. Since H' has no automorphisms itself, h must move all of its vertices. Furthermore, since H' has more vertices than G, h must send a vertex of H' to another vertex of H'. This means that as a permutation of the vertices of the component $H \cup H'$, h is completely determined: h must be a flip of $H \cup H'$ about some vertex of H'. This cannot happen, since by construction H' contains a cycle if and only if H does not.

Second, v_1 has larger degree in G' than in G, so there are no automorphisms of G' mapping v_1 to any other vertex v_2, \ldots, v_{p^m} . Since g maps v_1 to v_2 , g does not extend to any automorphism of G'.

Hence by induction G' has fixing number at most k - 1. If S is a fixing set of G' with $|S| \le k - 1$, then $S' = S \cup \{v_1\}$ is a fixing set of G with $|S'| \le k$. Therefore G has fixing number at most k, and fix $(\Gamma) = \{1, \ldots, k\}$.

3.2 Symmetric groups

The *inflation* of a graph G, Inf(G), is a graph with a vertex for each ordered pair (v, e), where v and e are a vertex and an edge of G, and v and e are incident. Inf(G) has an edge between (v_1, e_1) and (v_2, e_2) if $v_1 = v_2$ or $e_1 = e_2$. We denote the k-fold inflation of the graph G by $Inf^k(G)$.

For a positive integer n, let G_k be the graph with a vertex for each sequence (x_1, \ldots, x_{k+1}) of k+1 integers from the set $\{1, \ldots, n\}$ with x_1 different from the remaining integers in the sequence. Vertices $u = (u_1, \ldots, u_{k+1})$ and $v = (v_1, \ldots, v_{k+1})$ are adjacent if and only if there exists some index i such that $u_j = v_j$ for all j < i, $u_i \neq v_i$, and $u_j = v_i$ and $v_j = u_i$ for all j > i.

Lemma 15. The graphs G_k and $\text{Inf}^k(K_n)$ are isomorphic.

Proof. We define an isomorphism $\varphi : \operatorname{Inf}^k(K_n) \to G_k$ inductively. For the base case, note that $\operatorname{Inf}^0(K_n) \cong G_0 \cong K_n$. Now assume $\varphi' : \operatorname{Inf}^{k-1}(K_n) \to G_{k-1}$ is an isomorphism, and suppose that v is a vertex in $\operatorname{Inf}^k(K_n)$. By the definition of the inflation, v = (v', e'), where v' is a vertex in $\operatorname{Inf}^{k-1}(K_n)$ and e' is an edge in $\operatorname{Inf}^{k-1}(K_n)$. So $\varphi'(v') = (a_1, \ldots, a_k)$ and $e' = \{v', u'\}$ where $\varphi'(u') = (b_1, \ldots, b_k)$, for two vertices (a_1, \ldots, a_k) and (b_1, \ldots, b_k) in G_{k-1} . Since $v' \sim u'$, by the definition of G_{k-1} , there exists an index $1 \leq i \leq k$ such

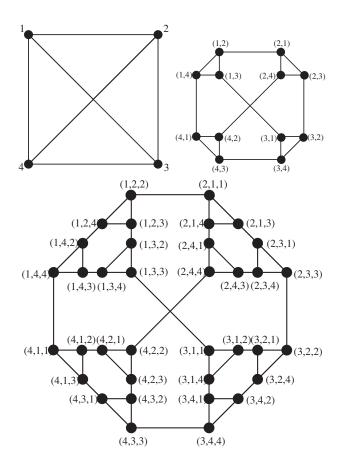


Figure 5: The graph K_4 and its first and second inflations.

that $a_j = b_j$ for all $1 \le j < i$, $a_i \ne b_i$, and $a_j = b_i$ and $b_j = a_i$ for all $i < j \le k$. We define $\varphi(v) = (a_1, \ldots, a_k, b_i)$. Note that since φ' is a bijection by induction, it is easy to see that φ is a bijection as well.

We now prove that φ is an isomorphism. First suppose that v and u are adjacent vertices of $\operatorname{Inf}^k(K_n)$. By the definition of inflation, v = (v', e') and u = (u', d') for two vertices v' and u' in $\operatorname{Inf}^{k-1}(K_n)$ and two edges e' and d' in $\operatorname{Inf}^{k-1}(K_n)$ incident to v' and u', respectively. By the definition of adjacency in $\operatorname{Inf}^k(K_n)$, either v' = u' or e' = d'.

Case 1. v' = u'. In this case, $\varphi'(v') = \varphi'(u') = (a_1, \ldots, a_k)$, so $\varphi(v)$ and $\varphi(u)$ differ only in their last coordinate. Therefore $\varphi(v) \sim \varphi(u)$ by the definition of adjacency in G_k .

Case 2. e' = d'. Since e' is incident to v' and d' is incident to u', e' = d' must be the edge between the vertices v' and u'. So $\varphi'(v') \sim \varphi'(u')$, hence $\varphi'(v')$ and $\varphi'(u')$ must satisfy the definition of adjacency in G_{k-1} . By the definition of φ , $\varphi(v)$ and $\varphi(u)$ are still adjacent in G_k .

Now suppose that v and u are non-adjacent vertices of $\operatorname{Inf}^k(K_n)$, and again let v = (v', e') and u = (u', d'). By the definition of adjacency in $\operatorname{Inf}^k(K_n)$, $v' \neq u'$ and $e' \neq d'$. **Case 1.** v' is not adjacent to u'. So $\varphi'(v') \not\sim \varphi'(u')$, so the sequences $\varphi'(v')$ and $\varphi'(u')$ do not satisfy the definition of adjacency in G_{k-1} . Since $\varphi(v)$ and $\varphi(u)$ are formed from $\varphi'(v')$ and $\varphi'(u')$ by appending an extra number to their sequences, the new sequences $\varphi(v)$ and $\varphi(u)$ still do not satisfy the definition of adjacency in G_k .

Case 2. v' is adjacent to u'. Since $v' \neq u'$, $\varphi'(v')$ and $\varphi'(u')$ differ in their *k*th coordinate. But since $e' \neq d'$, either the (k+1)st coordinate of $\varphi(v)$ differs from the *k*th coordinate of $\varphi(v)$, or the (k+1)st coordinate of $\varphi(u)$ differs from the *k*th coordinate of $\varphi(u)$. Therefore $\varphi(v)$ is not adjacent to $\varphi(u)$ in G_k .

By Lemma 15, we may label the vertices of $\text{Inf}^k(K_n)$ using the vertices of G_k , and follow the rule for adjacency of vertices in $\text{Inf}^k(K_n)$ given by the definition of G_k . We do this for the remainder of this section.

Theorem 16. For n > 3 and $k \ge 0$, $\operatorname{Aut}(\operatorname{Inf}^k(K_n)) = S_n$ and $\operatorname{fix}(\operatorname{Inf}^k(K_n)) = \lceil \frac{n-1}{k+1} \rceil$.

Proof. The statement is clear for k = 0, so assume k > 0. Since each vertex of $\text{Inf}^k(K_n)$ is labeled with a sequence of the numbers $\{1, \ldots, n\}$ of length k + 1 by Lemma 15, every permutation g in S_n induces a natural permutation of the vertices of $\text{Inf}^k(K_n)$. Again by Lemma 15, it is easy to see that these permutations are all automorphisms of $\text{Inf}^k(K_n)$. So $S_n < \text{Aut}(\text{Inf}^k(K_n))$.

Now suppose that $g \in \operatorname{Aut}(\operatorname{Inf}^k(K_n))$. We show that g is determined as a permutation of the numbers 1 through n in the labeling sequences of the vertices of $\operatorname{Inf}^k(K_n)$, and therefore $g \in S_n$. Suppose $v = (a_1, \ldots, a_{k+1})$ and $w = (b_1, \ldots, b_{k+1})$ are two vertices in $\operatorname{Inf}^k(K_n)$. By the definition of adjacency in G_k , if $a_i = b_i$ for $1 \leq i \leq k$, then v and w are adjacent. Therefore if we partition $\operatorname{Inf}^k(K_n)$ into blocks of vertices with the same first kelements in their labeling sequence, each block forms a maximal clique of $\operatorname{Inf}^k(K_n)$. The graph formed by contracting each of these maximal cliques to a single vertex is $\operatorname{Inf}^{k-1}(K_n)$. Since maximal cliques are preserved under automorphisms, the automorphism g induces a natural automorphism g' on $\operatorname{Inf}^{k-1}(K_n)$. By induction, g' is determined as a permutation p of the numbers 1 through n in the labeling sequences of the vertices of $\operatorname{Inf}^{k-1}(K_n)$. Now g is determined by the same permutation p, since the action of p on (a_1, \ldots, a_k) determines which maximal clique contains g(v), and the action of p on a_{k+1} determines g(v) within that maximal clique.

By the definition of the correspondence between an element g of $\operatorname{Aut}(\operatorname{Inf}^k(K_n))$ and its corresponding permutation p in S_n , for any vertex $v = (a_1, \ldots, a_{k+1})$ of $\operatorname{Inf}^k(K_n), g(v) = v$ if and only if $p(a_i) = a_i$ for all $1 \leq i \leq k+1$. Therefore $\operatorname{stab}(v) = \operatorname{stab}(\{a_1, \ldots, a_{k+1}\})$. This means that any set of vertices whose vertex labels include the set $\{1, \ldots, n-1\}$ is a fixing set of $\operatorname{Inf}^k(K_n)$. One such set is $\{(1, \ldots, k+1), (k+2, \ldots, 2k+1), \ldots, (mk+$ $m+1, \ldots, mk+m+k+1), (n-k-1, \ldots, n-1)\}$, where $m = \lfloor \frac{n-1}{k+1} \rfloor$. This set has $\lceil \frac{n-1}{k+1} \rceil$ vertices. Conversely, any set S of vertices whose vertex labels do not include any two of the numbers 1 through n, say i and j, cannot be a fixing set, since the element of $\operatorname{Aut}(\operatorname{Inf}^k(K_n))$ corresponding to the transposition (i, j) is a non-identity element of the stabilizer of S. This clearly requires at least $\lceil \frac{n-1}{k+1} \rceil$ vertices, so fix $(\operatorname{Inf}^k(K_n)) = \lceil \frac{n-1}{k+1} \rceil$. \Box

It seems likely that the proof of Theorem 16 could extend to inflations of graphs other than K_n . However, since $\text{Inf}^k(C_n) = C_{2^k n}$, $\text{fix}(\text{Inf}^k(C_n)) = 2$ for all $k \ge 0$ and $n \ge 3$. This motivates the following question.

Open Question. For which graphs G is it true that $fix(Inf^k(G)) = \lceil \frac{fix(G)}{k+1} \rceil$?

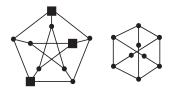


Figure 6: The Petersen graph with a fixing set shown as square vertices, and the Petersen graph with one vertex deleted.

Proposition 17. The Petersen graph P has automorphism group S_5 and fixing number 3.

Proof. Many proofs that $\operatorname{Aut}(P) = S_5$ appear in the literature; one can be found in [2]. A fixing set of P with 3 vertices is shown in Figure 6. It remains to show that any fixing set of P has at least 3 vertices. Suppose that $S = \{v_1, \ldots, v_k\}$ is a fixing set of P. Since P is vertex-transitive [2], we may choose v_1 to be any vertex of P. Since automorphisms in $\operatorname{stab}(v_1)$ preserve distance from v_1 , any element of $\operatorname{stab}(v_1)$ must permute the three vertices adjacent to v_1 among themselves, and the six vertices that are distance two from v_1 among themselves. Since automorphisms of $P - v_1$ also have this property, fixing the rest of P is equivalent to fixing the graph $P - v_1$. This graph is shown in Figure 6, and has fixing number 2 since its automorphisms are the same as the automorphisms of C_6 .

Lemma 18. For any positive integer n, if i is a prime power dividing n!, and j is the number of prime factors of n!/i, counting multiplicities, then $\max(\operatorname{fix}(S_n)) \leq j + 1$.

Proof. Let G be a graph with $\operatorname{Aut}(G) = S_n$. Let g be an element of S_n with order i. Since i is a prime power, by Lemma 11, as a permutation of the vertices of G, g contains a cycle of order i. Let v be a vertex in this cycle, and fix v. Since g is not an element of $\operatorname{stab}(v)$, $|\operatorname{stab}(v)| \leq n!/i$. Hence G can be fixed with j additional vertices by Lemma 7.

We conjecture that this lemma can be improved by fixing more than one vertex. However, one cannot use induction since the group $\operatorname{stab}(v)$ in the proof of Lemma 18 may not be symmetric.

We also have an upper bound on $\max(\operatorname{fix}(S_n))$ given by the following lemma, which appears in [4].

Lemma 19. $l(S_n) = \lceil 3n/2 \rceil - b(n) - 1$, where b(n) is the number of ones in the binary representation of n.

The following table gives lower and upper bounds on the set $fix(S_n)$, given by Propositions 6, 16, 17, 18, and 19. Note that Lemma 18 is the better upper bound for $n \leq 8$, and Lemmas 6 and 19 are better for $n \geq 10$.

group	lower bound	upper bound
S_2	$\{1\}$	{1}
S_3	$\{1,2\}$	$\{1,2\}$
S_4	$\{1,2,3\}$	$\{1,2,3\}$
S_5	$\{1,2,3,4\}$	$\{1,2,3,4\}$
S_6	$\{1,2,3,5\}$	$\{1,2,3,4,5,6\}$
S_7	$\{1,2,3,6\}$	$\{1,2,3,4,5,6,7\}$
S_8	$\{1,2,3,4,7\}$	$\{1,2,3,4,5,6,7,8,9\}$
S_9	$\{1,2,3,4,8\}$	$\{1,2,3,4,5,6,7,8,9,10,11\}$
S_{10}	$\{1,2,3,5,9\}$	$\{1,2,3,4,5,6,7,8,9,10,11,12\}$

Motivated by the first four rows of the table, we make the following conjecture.

Conjecture 20. $fix(S_n) = \{1, ..., n-1\}.$

Of particular interest is the potential gap which occurs first in fix (S_6) . More generally, all known examples of fixing sets of non-trivial finite groups are of the form $\{1, \ldots, k\}$ for some k. If the fixing set of every non-trivial finite group is of this form, then the computation of a fixing set becomes much easier: we need only to find the largest value in the set, which we may then call the *fixing number* of the group.

Open Question. For every non-trivial finite group Γ , does there exist a positive integer k such that fix $(\Gamma) = \{1, \ldots, k\}$?

Acknowledgements

We thank Pete L. Clark and an anonymous reviewer for many helpful suggestions.

References

- Michael O. Albertson and Karen L. Collins. Symmetry breaking in graphs. *Electron. J. Combin.*, 3(1):Research Paper 18, approx. 17 pp. (electronic), 1996.
- [2] Lowell W. Beineke and Robin J. Wilson, editors. Graph connections: Relationships between graph theory and other areas of mathematics, volume 5 of Oxford Lecture Series in Mathematics and its Applications. The Clarendon Press Oxford University Press, New York, 1997.
- [3] Debra Boutin. Identifying graph automorphisms using determining sets. *Electron. J. Combin.*, 13(1):Research Paper 78, approx. 14 pp. (electronic), 2006.
- [4] Peter J. Cameron, Ron Solomon, and Alexandre Turull. Chains of subgroups in symmetric groups. J. Algebra, 127(2):340–352, 1989.
- [5] Karen Collins and Joshua D. Laison. Fixing numbers of Kneser graphs. preprint, 2008.

- [6] Karen L. Collins and Ann N. Trenk. The distinguishing chromatic number. *Electron. J. Combin.*, 13(1):Research Paper 16, 19 pp. (electronic), 2006.
- [7] John D. Dixon and Brian Mortimer. Permutation groups, volume 163 of Graduate Texts in Mathematics. Springer-Verlag, New York, 1996.
- [8] David S. Dummit and Richard M. Foote. Abstract algebra. John Wiley and Sons, Inc., Hoboken, NJ, 3 edition, 2004.
- [9] David Erwin and Frank Harary. Destroying automorphisms by fixing nodes. Discrete Math., 306(24):3244–3252, 2006.
- [10] Robert Frucht. Hertellung von graphen mit vorgegebenen abstrakten gruppen. Compositio Math., 6:239–250, 1938.
- [11] Robert Frucht. Graphs of degree three with a given abstract group. Canadian J. Math., 1:365–378, 1949.
- [12] Chris Godsil and Gordon Royle. Algebraic graph theory, volume 207 of Graduate Texts in Mathematics. Springer-Verlag, New York, 2001.
- [13] Julianna Tymoczko. Distinguishing numbers for graphs and groups. Electron. J. Combin., 11(1):Research Paper 63, 13 pp. (electronic), 2004.