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# A New Method for Nonsmooth Convex Optimization\*

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A new method for minimizing a proper closed convex function f is proposed and its convergence properties are studied. The convergence rate depends on both the growth speed of f at minimizers and the choice of proximal parameters. An application of the method extends the corresponding results given by Kort and Bertsekas for proximal minimization algorithms to the case in which the iteration points are calculated approximately. In particular, it relaxes the convergence conditions of Rockafellar's results for the proximal point algorithm.

Keywords: Proper closed convex function; Proximal minimization; Global and superlinear convergence

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#### **1 INTRODUCTION**

Consider the proximal minimization algorithm for solving the following convex minimization problem

$$\min\{f(x): x \in \mathbb{R}^n\}$$
(1)

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defined by

$$x_{k+1} = \arg\min\left\{f(x) + \frac{1}{(\tau+1)\lambda_k} \|x - x_k\|^{\tau+1}: x \in \mathbb{R}^n\right\},\$$

where f is a proper closed convex function defined in  $\mathbb{R}^n$ , the proximal parameters  $\lambda_k$  and  $\tau$  are positive constants,  $\|\cdot\|$  denotes the standard Euclidean norm on  $\mathbb{R}^n$ , (see Martinet [15] which is the case  $\tau = 1$  based on [16]). The algorithm is globally convergent, and its local speed of convergence depends on: (i) The growth speed of f at the minimizers; (ii) The choice of  $\lambda_k$ ; and (iii) The choice of  $\tau$  (see [1,2,11]). For a survey of results of the proximal type minimization we refer the reader to [3-6,8-10,14,20-24], especially to [3,5,9,20].

Rockafellar [20] made a nice extension for the proximal point algorithm. In his algorithm, the iteration points are calculated approximately. A question is: do corresponding results of Kort and Bertsekas [11] hold for the algorithm of Rockafellar [20] when it is applied to a proper closed convex function?

In the next section, a new method for solving nonsmooth convex optimization problems is presented. This algorithm is globally convergent, and its local speed of convergence depends on both the growth speed at the minimizers and the choice of proximal parameters. These results are similar to those in [1,2,11]. Furthermore, a general proximal point algorithm introduced in [20] can be regarded as a special case of this method. In Section 3, an application of this algorithm relaxes a key condition of Rockafellar [20] for convergence of the proximal point algorithm. The condition,

$$\sum_{k=1}^{\infty} \delta_k < +\infty,$$

becomes

$$\limsup\{\delta_k\}<1,$$

where  $\delta_k$  is a parameter in the algorithm. In the same section, some special cases of the algorithm are also discussed. Some concluding remarks are in Section 4.

#### 2 METHOD AND ITS CONVERGENCE PROPERTIES

## 2.1 The Algorithm and its Global Convergence

Algorithm 1 Let q > 1 be a given positive number. Let  $x_1 \in \mathbb{R}^n$  be arbitrary. In general, given  $x_k \in \mathbb{R}^n$ , generate a proximal parameter  $t_k > 0$  and  $x_{k+1} \in \mathbb{R}^n$  such that there exists  $g_{k+1} \in \partial f(x_{k+1})$  satisfying the following inequality:

$$g_{k+1}^T(x_k - x_{k+1}) \ge t_k \|g_{k+1}\|^q.$$
(2)

THEOREM 2.1 Let  $\{x_k\}_{k=1}^{\infty}$  be any sequence generated by Algorithm 1.

- (i) If ∑<sub>k=1</sub><sup>∞</sup> t<sub>k</sub> = +∞, then either f(x<sub>k</sub>) → -∞ or lim inf<sub>k→∞</sub> ||g<sub>k</sub>|| = 0. In the latter case, if {x<sub>k</sub>}<sub>k=1</sub><sup>∞</sup> is a bounded set, then every accumulation point of {x<sub>k</sub>}<sub>k=1</sub><sup>∞</sup> is an optimal point of (1).
- (ii) If inf{t<sub>k</sub>} > 0, then either f(x<sub>k</sub>) → -∞ or g<sub>k</sub> → 0. In the latter case, every accumulation point of {x<sub>k</sub>}<sup>∞</sup><sub>k=1</sub> (if there is any) is an optimal point of (1).

Furthermore, Let  $X^*$  be the set of minimizers of f,  $f^* = \inf\{f(x): x \in \mathbb{R}^n\}$ . Then

$$f(x_{k+1}) - f^* \le \inf\{\|x_{k+1} - x^*\|: x^* \in X^*\}[t_k^{-1}\|x_{k+1} - x_k\|]^{(q-1)^{-1}}$$

*Proof* From the subgradient inequality,

$$f(x_k) \ge f(x_{k+1}) + g_{k+1}^T(x_k - x_{k+1}).$$

By (2), we have

$$f(x_k) \ge f(x_{k+1}) + t_k \|g_{k+1}\|^q.$$
(3)

This implies that  $f(x_{k+1}) \leq f(x_k)$ . If the decreasing sequence  $\{f(x_k)\}_{k=1}^{\infty}$  tends to  $-\infty$ , we are done. Otherwise, this sequence is bounded from below.

By (3),

$$\sum_{k=1}^{\infty} t_k \|g_{k+1}\|^q < +\infty.$$

If  $||g_k||$  is bounded below by a positive number, then the above inequality implies that  $\sum_{k=1}^{\infty} t_k < +\infty$ , which contradicts our assumptions in (i). Thus,  $\lim \inf_{k\to\infty} ||g_k|| = 0$ . If  $\{x_k\}_{k=1}^{\infty}$  is a bounded set, then there exists an infinite index set K such that  $\lim_{k \in K} g_k = 0$ . Without loss of generality, we may assume that  $\lim_{k \in K} x_k = x^*$ .

Applying

$$f(x) \ge f(x_k) + g_k^T(x - x_k)$$

and letting  $k \to \infty$ ,  $k \in K$ , we have that for all  $x \in \mathbb{R}^n$ ,  $f(x) \ge f(x^*)$ , which implies that  $x^*$  is a minimum point of f. By  $f(x_{k+1}) \le f(x_k)$ , we deduce that every accumulation point of  $\{x_k\}_{k=1}^{\infty}$  has the same optimal value,  $f(x^*)$ , and is, hence, also a minimum point of (1).

We now prove the second conclusion. From (3), we have that either  $f(x_k) \to -\infty$  or  $g_k \to 0$  by using  $\inf\{t_k\} > 0$ . Let  $\{x_k: k \in K\}$  be any convergence subsequence of  $\{x_k\}_{k=1}^{\infty}$  and  $\lim_{k \in K} x_k = x^*$ . Then,

$$\lim_{k \in K} \partial f(x_k) \subseteq \partial f(x^*)$$

yields that  $0 \in \partial f(x^*)$ . So  $x^*$  is an optimal point of (1), since f is a convex function.

For each  $x^* \in X^*$  and any k, since

$$f^* = f(x^*) \ge f(x_{k+1}) + g_{k+1}^T (x^* - x_{k+1}),$$

by (2), we have

$$f(x_{k+1}) - f^* \le ||g_{k+1}|| ||x_{k+1} - x^*||$$
  
$$\le ||x_{k+1} - x^*||[t_k^{-1}||x_{k+1} - x_k||]^{(q-1)^{-1}}$$

Hence

$$f(x_{k+1}) - f^* \le \inf\{\|x_{k+1} - x^*\|: x^* \in X^*\}[t_k^{-1}\|x_{k+1} - x_k\|]^{(q-1)^{-1}}.$$

The second conclusion follows.

COROLLARY 2.1 Suppose that  $X^*$  is nonempty, compact and  $\{t_k\}_{k=1}^{\infty}$  is bounded away from zero. Let  $\{x_k\}_{k=1}^{\infty}$  be any sequence generated by

Algorithm 1. Then  $\{x_k\}_{k=1}^{\infty}$  is bounded,  $f(x_k) \rightarrow f^*$  and

$$\lim_{k \to \infty} \rho(x_k; X^*) = 0, \tag{4}$$

where  $\rho(x; X^*)$  is the distance from x to  $X^*$  given by

$$\rho(x; X^*) = \min\{\|x - x^*\|: x^* \in X^*\}.$$

**Proof** By Theorems 8.4 and 8.7 of [18], the level sets of f are bounded. On the other hand, for all k,  $f(x_{k+1}) \ge f(x_k)$ . Hence  $\{x_k\}_{k=1}^{\infty}$  is contained in some level sets of f. Therefore,  $\{x_k\}_{k=1}^{\infty}$  is a bounded set. By (ii) of Theorem 2.1 and the compactness of  $X^*$ , we have  $f(x_k) \to f^*$  and (4).

#### 2.2 Local Convergence

We from now on assume that  $\inf\{t_k\} > 0$ . Like [3,11], we need the following key assumption which was used to analyze the convergence rate of the proximal minimization algorithm for quadratic as well as certain types of nonquadratic proximal terms (see [3,11]).

Assumption 1  $X^*$  is nonempty and compact. Furthermore, there exist scalars  $\alpha > 0$ ,  $p \ge 1$  and  $\delta > 0$  such that for any x with  $\rho(x; X^*) \le \delta$ ,

$$f(x) - f^* \ge \alpha(\rho(x; X^*))^p.$$
<sup>(5)</sup>

LEMMA 2.1 Suppose that (5) holds. Then for each x with  $\rho(x; X^*) \leq \delta$ and  $g \in \partial f(x)$ ,

$$\|g\| \ge \alpha^{1/p} [f(x) - f^*]^{(p-1)/p}.$$
(6)

*Proof* Since for any  $g \in \partial f(x)$  and any  $x^* \in X$ ,

$$f^* = f(x^*) \ge f(x) + g^T(x^* - x).$$

So we have

$$f^* - f(x) \ge \max\{g^T(x^* - x): x^* \in X^*\} \\ \ge \max\{-\|g\| \|x^* - x\|: x^* \in X^*\} \\ = -\|g\|\rho(x; X^*).$$

Hence,

$$f(x) - f^* \le ||g|| \rho(x; X^*) \le ||g|| \left[\frac{f(x) - f^*}{\alpha}\right]^{1/p}$$

by using (5). This implies (6).

LEMMA 2.2 Suppose that (5) holds,  $\{x_k\}_{k=1}^{\infty}$  is generated by Algorithm 1. Then for all large k,

$$\frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} \le \frac{1}{1 + t_k \alpha^{q/p} [f(x_{k+1}) - f^*]^{((q-1)(p-1)-1)/p}}.$$
 (7)

*Proof* Since  $\rho(x_k; X^*) \to 0$  and  $f(x_k) > f(x_{k+1})$ , we have  $k_1 > 1$  such that for all  $k \ge k_1$ ,  $\rho(x_k; X^*) \le \delta$ . Using  $g_{k+1} \in \partial f(x_{k+1})$ , we have

$$f(x_k) \ge f(x_{k+1}) + g_{k+1}^T(x_k - x_{k+1})$$
  

$$\ge f(x_{k+1}) + t_k ||g_{k+1}||^q$$
  

$$\ge f(x_{k+1}) + t_k \alpha^{q/p} [f(x_{k+1}) - f^*]^{q(p-1)/p}$$

by Lemma 2.1. Therefore,

$$f(x_k) - f^* \ge f(x_{k+1}) - f^* + t_k \alpha^{q/p} [f(x_{k+1}) - f^*]^{q(p-1)/p}.$$

This implies the conclusion by using  $f(x_{k+1}) > f^*$ .

THEOREM 2.2 Suppose that (5) holds,  $\{x_k\}_{k=1}^{\infty}$  is generated by Algorithm 1. Then  $f(x_k)$  tends to  $f^*$  superlinearly under any one of the following conditions ((a), (b) or (c)):

- (a) (q-1)(p-1) < 1.
- (b) (q-1)(p-1) = 1 and  $t_k \to +\infty$ . (If  $t_k \to t_* < +\infty$ , then  $f(x_k)$  tends to  $f^*$  linearly.)
- (c) (q-1)(p-1) > 1 and

$$\lim_{k \to \infty} t_k [f(x_k) - f^*]^{((q-1)p-q)/p} = +\infty.$$
(8)

If p = 1, then there exists  $k_2 > k_1$  such that  $x_{k_2} \in X^*$ , i.e., the minimizer of f can be determined within a finite number of steps.

Furthermore, if there exists a scalar  $\beta < +\infty$  such that

$$f(x) - f^* \le \beta(\rho(x; X^*))^p \quad \text{for all } x \text{ with } \rho(x; X^*) \le \delta, \qquad (9)$$

then  $x_k$  tends to  $X^*$  superlinearly under any of the above conditions ((a), (b) or (c)).

*Proof* The fact,  $f(x_{k+1}) \rightarrow f^*$ , implies that

$$\lim_{k \to \infty} t_k \alpha^{q/p} [f(x_{k+1}) - f^*]^{((q-1)(p-1)-1)/p} = +\infty,$$
(10)

if one of the conditions (a)-(b) holds. Then

$$\lim_{k \to \infty} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} = 0$$
(11)

by Lemma 2.2.

In case (b) with  $t_k \rightarrow t_* < +\infty$ , from (7), we have

$$\lim_{k \to \infty} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} \le \frac{1}{1 + t_* \alpha^{q/p}} < 1.$$

Hence  $f(x_k)$  tends to  $f^*$  linearly.

For (c), using

$$\frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} \le 1 \Big/ \Big\{ 1 + \alpha^{q/p} [f(x_{k+1}) - f^*/f(x_k) - f^*]^{((q-1)(p-1)-1)/p} \\ t_k [f(x_k) - f^*]^{((q-1)(p-1)-1)/p} \Big\},$$

(11) follows.

Recalling (5) and  $\rho(x_k; X^*) \le \delta$  for all  $k > k_1$ , which has been used in the proof of Lemma 2.2, we have

$$f(x_k) - f^* \ge \alpha (\rho(x_k; X^*))^p.$$
(12)

But from the proof of the Lemma 2.1,

$$f(x_k) - f^* \le \|g_k\| \rho(x_k; X^*).$$
(13)

Thus,

$$\alpha(\rho(x_k;X^*))^p \leq \|g_k\|\rho(x_k;X^*).$$

If p = 1, since  $\lim_{k \to \infty} g_k = 0$  and  $X^*$  is a compact set, we have  $k_2 > k_1$  such that  $x_{k_2} \in X^*$  (otherwise  $\alpha \le 0$ ).

We now prove the second conclusion. Let  $||x_k - X^*||$  denote  $\rho(x_k; X^*)$ . By (5) and (9) we have  $k_3 > k_1$  such that for all  $k > k_3$ ,

$$\alpha \|x_k - X^*\|^p \le f(x_k) - f^* \le \beta \|x_k - X^*\|^p.$$
(14)

This yields that  $f(x_k)$  tends to  $f^*$  superlinearly if and only if  $x_k$  tends to  $X^*$  superlinearly. The second conclusion follows from (11).

COROLLARY 2.2 Suppose that  $X^*$  is singleton, i.e.,  $X^* = \{x^*\}, \{x_k\}_{k=1}^{\infty}$  is generated by Algorithm 1.

(i) Assume that (5) and (9) hold. If there exist  $\tau_0 > \max\{0, (q-1)p-q\}$  and  $\tau_1 > 0$ , such that for all k,

$$t_k \|x_k - x_{k-1}\|^{\tau} \geq \tau_1,$$

then

$$\liminf_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0;$$

if, for all k,

$$t_k \|x_{k+1} - x_k\|^{\tau} \ge \tau_1,$$

then

$$\lim_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0.$$

(ii) Assume that (5) holds. If there exists  $\tau_{1'} > 0$  such that for all k,

$$t_k [f(x_k) - f(x_{k-1})]^{q-1} \ge \tau_{1'},$$

then

$$\liminf_{k \to \infty} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} = 0;$$

if, for all k,

$$t_k [f(x_{k+1}) - f(x_k)]^{q-1} \ge \tau_{1'},$$

then

$$\lim_{k \to \infty} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} = 0.$$

Furthermore, if (9) is true, then

$$\liminf_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0$$

or

$$\lim_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0$$

holds respectively.

*Proof* (i) It is very clear that  $t_k \to +\infty$ . Hence, we only need to consider case (c) of Theorem 2.2. If the conclusion is false, then we have a positive number  $\tau_2 < +\infty$ , such that for all k,

$$\frac{\|x_k - x^*\|}{\|x_{k-1} - x^*\|} \ge \frac{1}{\tau_2}.$$

Thus, we have

$$\frac{\|x_k - x^*\|}{\|x_k - x_{k-1}\|} \ge \frac{\|x_k - x^*\|}{\|x_k - x^*\| + \|x_{k-1} - x^*\|}$$
$$= \frac{1}{1 + \|x_{k-1} - x^*\| / \|x_k - x^*\|}$$
$$\ge \frac{1}{1 + \tau_2}.$$

This inequality and (5) yield that

$$t_{k}[f(x_{k}) - f^{*}]^{((q-1)p-q)/p}$$

$$= t_{k}||x_{k} - x_{k-1}||^{\tau} \left[\frac{||x_{k} - x^{*}||}{||x_{k} - x_{k-1}||}\right]^{[(q-1)p-q]} \left[\frac{f(x_{k}) - f^{*}}{||x_{k} - x^{*}||^{p}}\right]^{((q-1)p-q)/p}$$

$$\times \frac{1}{||x_{k} - x_{k-1}||^{\tau-[(q-1)p-q]}}$$

$$\geq \tau_{1} \alpha^{((q-1)p-q)/p} \left[\frac{1}{1+\tau_{2}}\right]^{[(q-1)p-q]} \frac{1}{||x_{k} - x_{k-1}||^{\tau-[(q-1)p-q]}}.$$

This last result implies that (8) holds. Hence  $x_k$  tends to  $x^*$  superlinearly by Theorem 2.2 and

$$\frac{1}{\tau_2} \le \frac{\|x_k - x^*\|}{\|x_{k-1} - x^*\|} \to 0.$$

This is a contradiction. So the first conclusion of (i) follows.

We now prove the second conclusion of (i). If the conclusion is false, then we have a positive number  $\tau_2 < +\infty$  and K such that for all  $k \in K$ ,

$$\frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} \ge \frac{1}{\tau_2}.$$

Using this inequality, similarly to the proof of the first conclusion, we have for all  $k \in K$ ,

$$t_k [f(x_{k+1}) - f^*]^{((q-1)p-q)/p} \\ \ge \tau_1 \alpha^{((q-1)p-q)/p} \left[\frac{1}{1+\tau_2}\right]^{[(q-1)p-q]} \frac{1}{\|x_{k+1} - x_k\|^{\tau-[(q-1)p-q]}}.$$

This and Lemma 2.2 imply that

$$\lim_{k \in K} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*} = 0.$$

By (5) and (9), we have

$$\lim_{k \in K} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} \le \left[\frac{\alpha}{\beta}\right]^{1/p} \lim_{k \in K} \left[\frac{f(x_{k+1}) - f^*}{f(x_k) - f^*}\right]^{1/p} = 0.$$

So  $\tau_2 = +\infty$ . Again, this is a contradiction. This proves the second conclusion of (i).

(ii) Assume that (5) holds. We prove the first conclusion. Using

$$\frac{f(x_k) - f^*}{f(x_k) - f(x_{k-1})} \ge \frac{1}{1 + (f(x_{k-1}) - f^*)/(f(x_k) - f^*)}$$

and

$$t_{k}[f(x_{k}) - f^{*}]^{((q-1)p-q)/p}$$
  
=  $t_{k} \Big[ f(x_{k}) - f(x_{k-1}) \Big]^{q-1} \Big[ \frac{f(x_{k}) - f^{*}}{f(x_{k}) - f(x_{k-1})} \Big]^{((q-1)p-q)/p}$   
 $\times \Big[ f(x_{k}) - f(x_{k-1}) \Big]^{-q/p},$ 

similarly to the proof of (i), if the conclusion is false, we have

$$\lim_{k \to \infty} t_k [f(x_k) - f^*]^{((q-1)p-q)/p} = +\infty.$$

Using conclusion (c) of Theorem 2.2, we have  $f(x_k)$  tends to  $f^*$  superlinearly which contradicts our assumption that the conclusion is false. Hence, the conclusion follows.

Similarly to the proof of the first conclusion of (ii) and the proof of the second conclusion of (i), we can prove the second conclusion of (ii).

If (9) holds, by

$$\frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} \le \left[\frac{\beta}{\alpha} \frac{f(x_{k+1}) - f^*}{f(x_k) - f^*}\right]^{1/p}$$

and the conclusions which we have just proved, the last conclusion of the corollary follows.

Using the results of this subsection, we may give some special cases of Algorithm 1 according to the choice of  $t_k$ . One such special case follows.

**Algorithm 1.1** In Algorithm 1, generate  $x_{k+1} \in \partial f(x_{k+1})$  satisfying the following inequality:

$$g_{k+1}^T(x_k - x_{k+1}) \ge ||x_k - x_{k-1}||^{-\tau} ||g_{k+1}||^q$$

or

$$g_{k+1}^T(x_k - x_{k+1}) \ge \left[f(x_k) - f(x_{k-1})\right]^{1-q} \|g_{k+1}\|^q.$$

In general, we can choose q = 2.

It is worth to note that, if we choose  $t_k$  as following:

$$t_k = r \, rac{g_k^T(x_{k-1} - x_k)}{\|g_k\|^q},$$

where r > 1 is a constant, then from (2), we can deduce that for all k,  $t_k \ge r^k$ . Hence  $t_k \to +\infty$ . Hence if p = q = 2, then, by Theorem 2.2,  $f(x_k)$  tends to  $f^*$  superlinearly under the condition (5).

Kort and Bertsekas in their novel paper [11] presented a combined primal-dual and penalty method for solving constrained convex programming. It essentially is a proximal point algorithm (see [2,19]). Hence, from the results of [1,2,11], under (5), the convergence rate of the ordinary proximal minimization algorithm is shown to depend on the growth speed of f at the minimizers as well as the proximal parameters  $\lambda_k$  and  $\tau$ . More precisely, if  $\tau = 1$ , the ordinary minimization algorithm has finite, superlinear and linear convergence rate depending on whether p=1, 1 and <math>p=2, respectively. The convergence rate is also superlinear if p=2 and  $\lambda_k \to \infty$ . In the case where the proximal term has a growth rate  $\tau > 1$  other than quadratic  $(\tau = 1)$ , the convergence rate is superlinear if 1 even in the $case where <math>p \ge 2$  (see [1,2,11]). From Theorem 2.2 (in the case of  $q = 1 + 1/\tau$ ), we obtain corresponding results for our algorithm.

## **3 APPLICATION TO THE PROXIMAL POINT METHOD**

In [20], Rockafellar presented two general proximal point algorithms for solving (1). One of them is as follows.

Algorithm A For  $x_k$ , generate  $x_{k+1} \in \mathbb{R}^n$  satisfying

$$\rho(0; S_k(x_{k+1})) \le \frac{\delta_k}{c_k} \|x_{k+1} - x_k\|,$$
(15)

where  $c_k > 0$ ,

$$\sum_{k=1}^{\infty} \delta_k < \infty, \tag{16}$$

and

$$S_k(x) = \partial f(x) + \frac{1}{c_k}(x - x_k).$$
(17)

In the following, we present a generalization of Algorithm A. In the new algorithm, we relax (16) to the following condition:

$$\limsup\{\delta_k\} < 1. \tag{18}$$

The new algorithm possesses a convergence rate similar to the rate for Algorithm A. The relaxation may be important in practice.

**Algorithm B** For any given  $\tau > 0$ , let  $x_1 \in \mathbb{R}^n$ . In general, for  $x_k \in \mathbb{R}^n$ , generate  $x_{k+1} \in \mathbb{R}^n$  satisfying

$$\rho(0; S_k^{\tau}(x_{k+1})) \le \frac{\delta_k}{c_k} \|x_{k+1} - x_k\|^{\tau},$$
(19)

$$S_k^{\tau}(x) = \partial f(x) + \frac{1}{c_k} \|x - x_k\|^{\tau - 1} (x - x_k).$$
<sup>(20)</sup>

The following theorem proves that Algorithm B is a special case of Algorithm 1, hence the conclusions of Theorem 2.2 hold for Algorithm B. It is very clear that the ordinary proximal minimization described in the introduction of this paper is a special case of Algorithm B. Therefore,

we extended the corresponding results [1,2,11], for ordinary proximal minimization algorithms to the case when the iteration points are calculated approximately.

**THEOREM 3.1** Suppose that  $x_{k+1}$  is generated by Algorithm B, and  $g_{k+1}$  is a vector in  $\partial f(x_{k+1})$  such that the norm of  $g_{k+1} + 1/c_k ||x_{k+1} - x_k||^{\tau-1}$   $(x_{k+1} - x_k)$  is not bigger than the right-hand side of (19). Then for all large k,  $(x_{k+1}, g_{k+1})$  satisfies (2) with  $q = 1 + 1/\tau$  and

$$t_k = \frac{1 - \delta_k}{(1 + \delta_k)^{1 + 1/\tau}} c_k^{1/\tau}.$$

*Proof* Without loss of generality, we may assume that

$$\sup\{\delta_k\} < 1. \tag{21}$$

By (19), (21) and (20), we have

$$\left\|g_{k+1} + \frac{1}{c_k}\|x_{k+1} - x_k\|^{\tau-1}(x_{k+1} - x_k)\right\| \le \frac{\delta_k}{c_k}\|x_{k+1} - x_k\|^{\tau}.$$
 (22)

The inequality,

$$\begin{split} \left[g_{k+1} + \frac{1}{c_k} \|x_{k+1} - x_k\|^{\tau-1} (x_{k+1} - x_k)\right]^T (x_{k+1} - x_k) \\ & \leq \left\|g_{k+1} + \frac{1}{c_k} \|x_{k+1} - x_k\|^{\tau-1} (x_{k+1} - x_k)\right\| \|x_{k+1} - x_k\|, \end{split}$$

and (22) imply that

$$g_{k+1}^T(x_{k+1}-x_k) \leq -\frac{1-\delta_k}{c_k} \|x_{k+1}-x_k\|^{\tau+1}.$$

Therefore,

$$g_{k+1}^{T}(x_{k} - x_{k+1}) \ge \frac{1 - \delta_{k}}{c_{k}} \|x_{k+1} - x_{k}\|^{\tau+1}.$$
 (23)

On the other hand, using (22), we have

$$\|g_{k+1}\| \le \frac{1+\delta_k}{c_k} \|x_{k+1} - x_k\|^{\tau}.$$
(24)

Now, the conclusion of the theorem follows (23) and (24).

By the above results, we may state global convergence and convergence rate results for Algorithm B. Noticing that  $t_k \rightarrow +\infty$  if and only if  $c_k \rightarrow +\infty$ , we only need to change q to  $1 + 1/\tau$  in Theorem 2.2 and Corollary 2.2. We do not go into this direction in details. We are more interested in the special case in which  $\tau = 1$ . In the following analysis, we always assume  $\tau = 1$ .

**THEOREM 3.2** Let  $\{x_k\}_{k=1}^{\infty}$  be any sequence generated by Algorithm B with  $\{c_k\}_{k=1}^{\infty}$  bounded away from zero. Suppose that  $\{x_k\}_{k=1}^{\infty}$  is bounded. Then every accumulation point of  $\{x_k\}_{k=1}^{\infty}$  is an optimal solution of (1). Furthermore, if  $x^*$  is an optimal solution of (1), then

$$f(x_{k+1}) - f(x^*) \le ||x_{k+1} - x^*|| \\ \times \left[ ||g_{k+1} + c_k^{-1}(x_{k+1} - x_k)|| + c_k^{-1} ||x_{k+1} - x_k|| \right] \\ \le \frac{1 + \delta_k}{c_k} ||x_{k+1} - x^*|| ||x_{k+1} - x_k||.$$

*Proof* By Theorem 2.1, we have the first conclusion. The proof of the first inequality in the second conclusion is similar to the proof of Theorem 4 in [18]. The second inequality follows (22).

For rates of convergence in optimization methods, a key condition is called the quadratic growth at  $x^*$  of f, i.e., in the neighbourhood of  $x^*$ ,  $O(x^*; r_1) = \{x: ||x-x^*|| < r_1\},\$ 

$$f(x) \ge f(x^*) + \alpha ||x - x^*||^2$$
 for all  $x \in O(x^*; r_1)$ , (25)

where  $\alpha > 0$ . Under (25) and some other additional assumptions  $(c_k \to +\infty \text{ and } \sum_{k=1}^{+\infty} \delta_k < +\infty)$ , Rockafellar proved the superlinear convergence of Algorithm A. The conditions in (16) and  $c_k \to +\infty$  are, however, quite strict. Very large  $c_k$  and very small  $\delta_k$ , in fact, imply that the calculation of  $x_{k+1}$  can be almost as difficult as the original minimization problem  $(0 \in \partial f(x_{k+1}))$ . Moreover, the condition (25) is

false for a large class convex functions (for example,  $f(x) = ||x||^p$ , p > 2). So it is necessary to discuss the possibility of relaxation of the requirements for  $c_k$ ,  $\delta_k$  and (25).

For given  $x \in \mathbb{R}^n$  and  $S \subseteq \mathbb{R}^+ = (0, +\infty)$ , a collection of proper closed convex functions (PCCF) is defined by

$$CC(x; S) = \{ f: f \in \text{PCCF and there exists } p \in S, 0 < \alpha_p(f; x) \\ \leq \beta_p(f; x) < \infty \},\$$

where

$$\alpha_p(f;x) = \lim_{y \to x} \inf \left\{ \frac{f(y) - f(x)}{\|y - x\|^p} \colon f(y) > f(x) \right\}$$
(26)

and

$$\beta_p(f;x) = \lim_{y \to x} \sup\left\{\frac{f(y) - f(x)}{\|y - x\|^p} \colon f(y) > f(x)\right\}.$$
 (27)

We now have the following theorem by Theorems 2.2 and 3.2.

THEOREM 3.3 Suppose that the conditions of Theorem 3.2 hold,  $\{x_k\}_{k=1}^{\infty}$  is generated by Algorithm B and  $x_k \rightarrow x^*$ . Then  $f(x_k)$  tends to  $f(x^*)$  superlinearly if one of the following conditions ((a<sub>1</sub>), (b<sub>1</sub>) or (c<sub>1</sub>)) holds:

- (a<sub>1</sub>)  $p \in (1, 2), \alpha_p(f; x^*) > 0;$
- (b<sub>1</sub>)  $p = 2, \alpha_2(f; x^*) > 0 \text{ and } c_k \to +\infty;$
- (c<sub>1</sub>)  $p \in (2, +\infty), \alpha_p(f; x^*) > 0$ , and  $\lim_{k \to \infty} c_k [f(x_k) f(x^*)]^{(p-2)/p} = +\infty$ .

If  $\alpha_1(f; x^*) > 0$ , then the minimizer of f can be determined after a finite number of steps.

Furthermore,  $x_k$  tends to  $x^*$  superlinearly if one of the following conditions ((a<sub>2</sub>) or (b<sub>2</sub>) or (c<sub>2</sub>)) holds:

(a<sub>2</sub>)  $f \in CC(x^*; (1, 2));$ (b<sub>2</sub>)  $f \in CC(x^*; \{2\})$  and  $c_k \to +\infty;$ (c<sub>2</sub>)  $f \in CC(x^*; (2, +\infty))$  and  $\lim_{k \to \infty} c_k [f(x_k) - f(x^*)]^{(p-2)/p} = +\infty.$  COROLLARY 3.1 Assume that f has a unique minimum, say  $x^*$ . Let  $\{x_k\}_{k=1}^{\infty}$  be a sequence generated by Algorithm B.

(i) If  $f \in CC(x^*; \{p\})$  and there exist  $\tau_3 > \max\{0, p-2\}$ , and  $\tau_4 > 0$ , such that for all k

$$c_k \|x_k - x_{k-1}\|^{\tau_3} \geq \tau_4,$$

then  $x_k$  tends to the unique minima of f, i.e.,  $x_k \rightarrow x^*$ , and

$$\liminf_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0;$$

if for all k,

$$c_k \|x_{k+1} - x_k\|^{\tau_3} \ge \tau_4,$$

then

$$\lim_{k\to\infty}\frac{\|x_{k+1}-x^*\|}{\|x_k-x^*\|}=0.$$

(ii) Assume  $\alpha_p(f; x^*) > 0$ . If there exists  $\tau_{1'} > 0$  such that for all k,

$$c_k[f(x_k)-f(x_{k-1})]\geq \tau_{4'},$$

then

$$\liminf_{k \to \infty} \frac{f(x_{k+1}) - f(x^*)}{f(x_k) - f(x^*)} = 0;$$

if, for all k,

$$c_k\big[f(x_{k+1})-f(x_k)\big]\geq \tau_{4'},$$

then

$$\lim_{k \to \infty} \frac{f(x_{k+1}) - f(x^*)}{f(x_k) - f(x^*)} = 0.$$

Furthermore, if  $f \in CC(x^*; p)$ , then

$$\liminf_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0$$

or

$$\lim_{k \to \infty} \frac{\|x_{k+1} - x^*\|}{\|x_k - x^*\|} = 0$$

holds respectively.

*Proof* Using the assumption of this theorem, we can deduce that  $\{x_k\}_{k=1}^{\infty}$  is bounded (see the proof of Corollary 2.1). By the boundedness of  $\{x_k\}_{k=1}^{\infty}$ ,

$$c_k \geq \frac{\tau_4}{\|x_k - x_{k-1}\|^{\tau_3}} \geq \frac{\tau_4}{2^{\tau_3} \sup\{\|x_k\|^{\tau_3}\}} > 0,$$

and Theorem 3.2, we have that  $x_k$  tends to  $x^*$ , the unique minimum of f, i.e.,  $x_k \rightarrow x^*$ . The conclusions follow Corollary 2.2 now.

It should be noted that the condition  $\alpha_2(f; x^*) > 0$  is weaker than (25). In fact,  $\alpha_2(f; x^*) > 0$  does not imply the uniqueness of the minimizers of f, but the following proposition holds.

**PROPOSITION 3.1** If (25) holds, then  $X^* = \{x^*\}$ .

Let  $Y^*$  be the set of all accumulation points of  $\{x_k\}_{k=1}^{\infty}$  generated by Algorithm B. Then we have the following result.

LEMMA 3.1 Suppose that  $Y^* \neq \emptyset$  has an isolated point and

$$\lim_{k \to \infty} \|x_{k+1} - x_k\| = 0.$$
 (28)

Then  $\{x_k\}_{k=1}^{\infty}$  is a convergent sequence.

*Proof* By the assumption on  $Y^*$ , there is a point  $y^* \in Y^*$  such that it is an isolation point in  $Y^*$ . Then there exist  $r_2 > 0$  and  $O(y^*; r_2) =$  $\{x: ||x - y^*|| < r_2\}$  such that  $Y^* \cap O(y^*; r_2) \setminus \{y^*\} = \emptyset$ . Let  $K(y^*) =$  $\{k: x_k \in O(y^*; r_2)\}$ . Then  $K(y^*)$  is an infinite set and

$$\lim_{k\in K(y^*)}x_k=y^*.$$

Furthermore, let  $K(y^*) = \{k_j: j_1 < j_2 \text{ implies } k_{j_1} \le k_{j_2}\}$  and N be the set of natural numbers. If  $N \setminus K(y^*)$  is an infinite set, then there exists an infinite set  $K_1(y^*) \subseteq K(y^*)$  such that for all  $k_i \in K_1(y^*)$ ,  $k_i + 1 \notin K(y^*)$ . Let  $l_1$  be the smallest number in  $N \setminus K(y^*)$ .

In general, suppose that we have  $\{l_1, l_2, ..., l_j\}$ . Then let  $l_{j+1}$  be the smallest number in  $N \setminus K(y^*) \setminus \{l_1, l_1 + 1, ..., l_j, l_j + 1\}$ . By induction, we have an infinite set  $\{l_i: i \in N\}$ . It is clear that  $K_1(y^*) = \{k_i = l_i - 1: i \in N\}$  is the desired set. Since  $\lim_{k_i \to \infty} x_{k_i} = y^*$ , we have  $\lim_{k_i \to \infty} x_{k_i+1} = y^*$  by (28). This is impossible from the definition of  $K(y^*)$  and  $k_i + 1 \notin K(y^*)$ . So  $N \setminus K(y^*)$  is a finite set. This implies that  $\lim_{k \to \infty} x_k = y^*$  by the definition of  $K(y^*)$ .

Indeed, if both  $L(f; x_1) = \{x: f(x) \le f(x_1)\}$  and  $\{c_k\}$  are bounded, then Theorem 3.2 implies that  $\lim_{k\to\infty} ||x_{k+1} - x_k|| = 0$ . In addition, if the condition (a<sub>2</sub>) in Theorem 3.3 holds and  $Y^*$  is not a dense subset of  $\mathbb{R}^n$ , then  $x_k$  tends to one of the minimizers of f superlinearly. In this sense, (25) relaxed by  $\alpha_2(f; x^*) > 0$  is a useful consideration for those problems with nonunique minima.

*Example 3.1* Let  $f(x) \in LEC$ , and  $f(x^*) = 0$ . For any given number l < -(p-2)/p, choose

$$c_{k} = \begin{cases} c > 0 & \text{if } 1 \le p < 2, \\ [f(x_{k})]^{l} & \text{if } p \ge 2 \end{cases}$$
(29)

in Algorithm B, then  $f(x_k)$  tends to  $f(x^*)$  superlinearly if  $\alpha_p(f; x^*) > 0$ . Furthermore,  $x_k$  tends to  $x^*$  superlinearly if  $f \in CC(x^*; \{p\})$ .

In Examples 3.2–3.4, we assume that  $c_k$  satisfies (29) in Algorithm B and that  $x_k$  tends to  $x^*$ .

*Example 3.2* Suppose that we wish to find a zero point of a function  $h(x) \in LEC$ . Then we can choose  $f(x) = \max\{0, h(x)\}$  and start Algorithm B from the point  $x_1$ ,  $h(x_1) > 0$ . Since  $f(x_k) > f(x_{k+1}) \rightarrow f(x^*)$ , also  $h(x_k) > h(x_{k+1}) \rightarrow h(x^*)$ . Thus,  $h(x^*) = 0$ . Furthermore, we have the following proposition.

**PROPOSITION 3.2** If  $\alpha_p(h; x^*) > 0$ , then  $h(x_k)$  tends to 0 superlinearly. If  $h \in CC(x^*; \{p\})$ , then  $x_k$  tends to  $x^*$  superlinearly.

*Proof* Since  $\alpha_p(f; x^*) = \alpha_p(h; x^*)$  and  $\beta_p(f; x^*) = \beta_p(h; x^*)$ , the conclusions follow from Theorem 3.3.

*Example 3.3* An important special case of Example 3.1 is called the inequality feasibility problem which is the most fundamental in convex mathematical programming.

Let  $f_i$  ( $i \in I$  a finite index set) be convex functions. Find  $x \in \mathbb{R}^n$  such that

$$f_i(x) \leq 0$$
 for all  $i \in I$ .

Let  $f(x) = \max\{0, f_i(x): i \in I\}$ . Then f is a convex function,  $f(x) \ge 0$ . Furthermore

$$f(x) = 0$$
 if and only if  $f_i(x) \le 0$  for all  $i \in I$ .

Hence, for the inequality feasibility problem, we have a method to solve it quickly if there exists  $p \in (0, +\infty)$  such that for all  $i \in I(x^*) =$  $\{i: f_i(x^*) = 0\}, f_i \in CC(x^*; \{p\})$ . More precisely, we have the following proposition.

**PROPOSITION 3.3** If for all  $i \in I(x^*)$ ,  $\alpha_p(f_i; x^*) > 0$ , then  $f(x_k)$  tends to 0 superlinearly; if for all  $i \in I(x^*)$ ,  $f_i \in CC(x^*; \{p\})$ , then  $x_k$  tends to  $x^*$  superlinearly.

**Proof** First of all, we note that  $I(x^*) \neq \emptyset$  from  $f(x_k) > 0$  and  $f(x_k) \rightarrow f(x^*) = 0$ . On the other hand, for  $i \notin I(x^*)$ , since  $f_i(x^*) < 0$ , so there exists  $r_2 > 0$ , such that for all  $x \in O(x^*; r_2)$ ,  $f_i(x) < 0$ . This implies that for  $i \notin I(x^*)$ ,

$$f_i(x) \neq f(x)$$
 for all  $x \in O(x^*; r_2)$ .

Hence, if for all  $i \in I(x^*)$ ,  $\alpha_p(f_i; x^*) > 0$ , then

$$\begin{aligned} \alpha_p(f;x^*) &= \lim_{x \to x^*} \inf \left\{ \frac{f(x) - f(x^*)}{\|x - x^*\|^p} \colon f(x) > f(x^*) \right\} \\ &\geq \lim_{x \to x^*} \inf \left\{ \frac{f_i(x) - f_i(x^*)}{\|x - x^*\|^p} \colon i \in I(x^*), f_i(x) = f(x) > 0 \right\} \\ &\geq \min\{\alpha_p(f_i;x^*) \colon i \in I(x^*)\} > 0. \end{aligned}$$

Therefore, the first conclusion follows Theorem 3.3.

By the same approach, if for all  $i \in I(x^*)$ ,  $f_i \in CC(x^*; \{p\})$ , we have

$$\beta_p(f; x^*) < +\infty.$$

The second conclusion follows Theorem 3.3.

*Example 3.4* Another case of Example 3.1 is finding the zero points of a given function  $F(x): \mathbb{R}^n \to \mathbb{R}^m$ . More precisely, find  $x^* \in \mathbb{R}^n$ , such that

$$F(x^*) = 0$$

if  $f(x) = \frac{1}{2}F(x)^T F(x)$  is a convex function (example: F(x) = Ax + b,  $A \in \mathbb{R}^{n \times m}$ ). For this problem, we have the following result.

**PROPOSITION 3.4** Suppose that f is a convex function from  $\mathbb{R}^n \to \mathbb{R}$ ,

$$\lim_{x \to x^*} \inf \left\{ \frac{\|F(x)\|^2}{\|x - x^*\|^p} : \ F(x) \neq 0 \right\} > 0.$$

Then  $||F(x_k)||$  tends to 0 superlinearly. Furthermore, if the following condition

$$\lim_{x \to x^*} \sup \left\{ \frac{\|F(x)\|^2}{\|x - x^*\|^p} : \ F(x) \neq 0 \right\} < +\infty$$

is added, then  $x_k$  tends to  $x^*$  superlinearly.

# 4 CONCLUDING REMARKS

The original motivation for the Moreau–Yosida regularization was to solve ill-conditioned systems of linear equations. It is noticed that Algorithm B always converges superlinearly when p < 2,  $c_k = constant$ ,  $\delta_k = \delta < 1$  and, for a large class of functions, a smaller p (<2) implies greater ill-conditioning. Comparing with other methods (for example, Newton or quasi-Newton methods), the proximal point method takes advantage of solving ill-conditioned problems (also see [13,17,7] for more details on this topic).

The benefit of Algorithm B can be explained by using the univariate function  $f(x) = \frac{1}{3}|x|^3$ . For this function, the algorithm given in [12] may

not converge superlinearly when starting from  $x_1 > 0$  (see Proposition 15 of [12]). If we choose  $c_k = [f(x_k)]^l$   $(l < -\frac{1}{3}$  is a constant) and apply Algorithm B, then  $\{x_k\}_{k=1}^{\infty}$  converges to the solution 0 superlinearly by the results of Example 3.1. Also, for this function, since (25) does not hold, the superlinear convergence results of [20] cannot be applied.

The results of Corollaries 2.2 and 3.1 are important for constructing a good algorithm in the spirit of the variable metric methods (see [12]). In [12], the authors gave a way to construct  $c_k$  such that  $c_k \to +\infty$ . They proved the algorithm possesses superlinear convergence under (25) when f is an  $LC^1$  function (in this case,  $f \in CC(x^*; \{2\})$ ). Since Algorithm 13 of [12] is a special case of Algorithm B (if for all  $k, \delta_k = 0$ ),  $x_k$  tends to  $x^*$  superlinearly from Theorem 3.3. On the other hand, if we let  $c_k = ||x_k - x_{k-1}||^{-\tau_3}$  in Algorithm B, then we can expect that  $x_k$ tends to  $x^*$  superlinearly even if f is not a smooth function using Corollary 3.1. Hence, the results of Corollary 3.1 provide a rule to construct the variable matrices based on the ideas of [12] (for penaltytype bundle algorithms) and quasi-Newton formulae.

It is not very hard to give a partial proximal minimization algorithm (see [3]) by a way similar to Algorithm 1. In fact, in Algorithm 1, if we let  $g_{k+1}$  be the vector with components zero for  $i \in I$  (where I is a subset of the index set  $\{1, \ldots, n\}$ , see [3]) then one of the special cases of the PPM algorithm given by [3] is contained in the Algorithm 1. The reader can go into this direction in details by following [3].

The essential content of this paper is a theoretical investigation of algorithms for nonsmooth convex optimization. We extended convergence results from those previously given but one question remains.

How can one solve (2) effectively if f is a nonsmooth function?

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